

# On the Need of Physical Security for Small Embedded Devices: A Case Study with COMP128-1 Implementations in SIM Cards

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Financial Cryptography and Data Security 2013

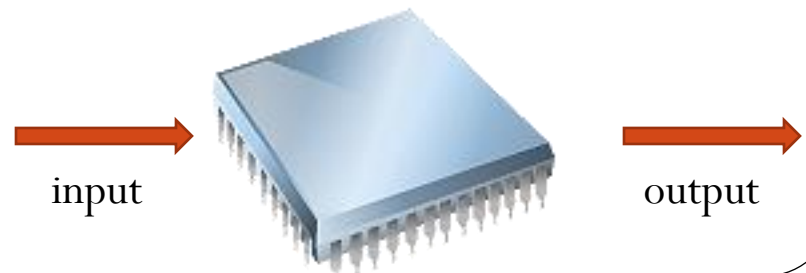
# Outline of the Talk

- Cryptography and Physical Security
- GSM and COMP128-1 (A3/A8) SIM cards
- Weakness and Attacks: Algorithmic vs. Physical
- A Case Study on COMP128-1 Implementations
- Lessons Learned

# How cryptography works?

- ▶ Typical Assumptions:
  - (1) A computational hard problem (RSA, AES ).
  - (2) Black-box: attacker **ONLY** sees input-output.
- ▶ Provable Security: Reductionist approach.

If one breaks the crypto-system (in polynomial-time), then it leads to efficient solution to the assumptions .
- ▶ Security guarantee **voided** if either assumption is not met.



# Are these assumptions safe?

- ▶ Typical Assumptions:

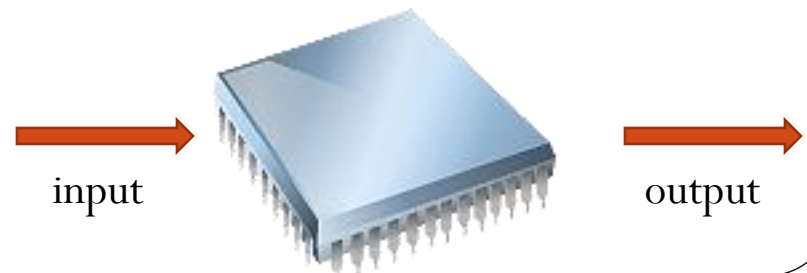
- (1) A computational hard problem (RSA, AES ).
- (2) Black-box: attacker **ONLY** sees input-output.

- ▶ Provable Security: Reductionist approach.

- ▶ Assumption #1 is ok (otherwise a breakthrough).

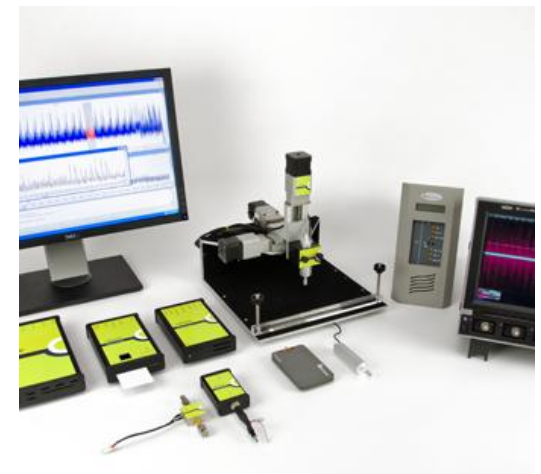
- ▶ Assumption #2 is not always respected.

The implementation of a cryptographic algorithm might be leaking in many forms.



# ● Side-channel attacks and beyond

- Definition: Any attack based on information gained from the physical implementation of a cryptosystem, rather than brute force or theoretical weaknesses in the algorithms.
- It takes many forms:
  - Timing Attacks
  - Power Analysis (PA)
  - Electro-Magnetic Analysis (EMA)
  - Acoustic Analysis
  - etc.
- More invasive physical attacks exist.



# ● Cryptographic Products in Real World

Smart cards equivalents, banking tokens, and other small embedded devices.



大卡、小卡对比

# Cellular networks (1-4G)

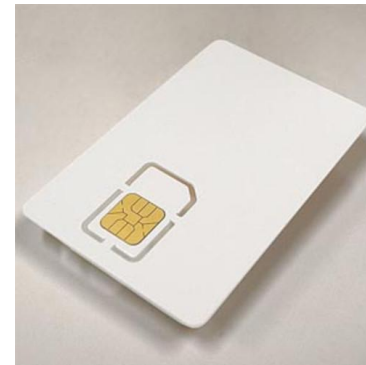
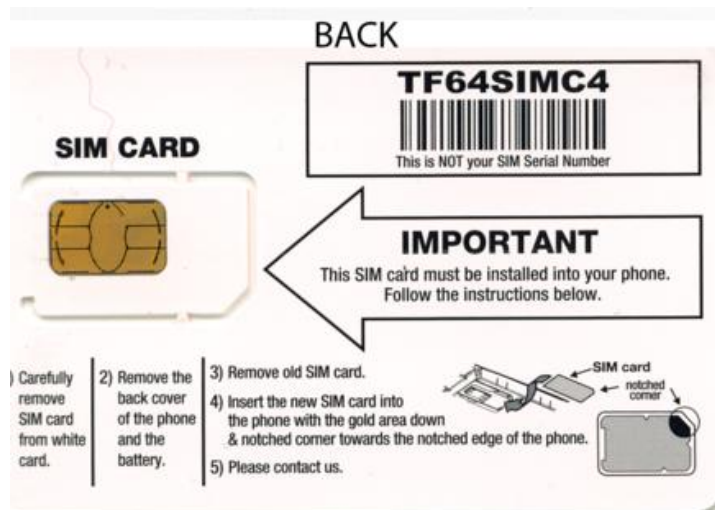
- 1G: analogue signal (last 90's)
- 2G: digital signal  
    **GSM** vs. CDMA
- 3G: UMTS vs. CDMA2000  
    high-speed data transmission
- 4G: LTE Advanced vs. WiMAX (IEEE 802.16e)



Despite the migration to 3G/4G, GSM remains the current dominant technology for mobile communications, especially in many developing countries.

# SIM cloning: the main threat to phone security

- SIM card is a smart card.
  - SIM stores: ICCID(serial number), IMSI (USER id), secret key K, contacts (optional).
- knowing IMSI and K allows one to clone the SIM card
- SIM Cloning : making fraudulent calls、 impersonation、 privacy breach、 internet banking security。

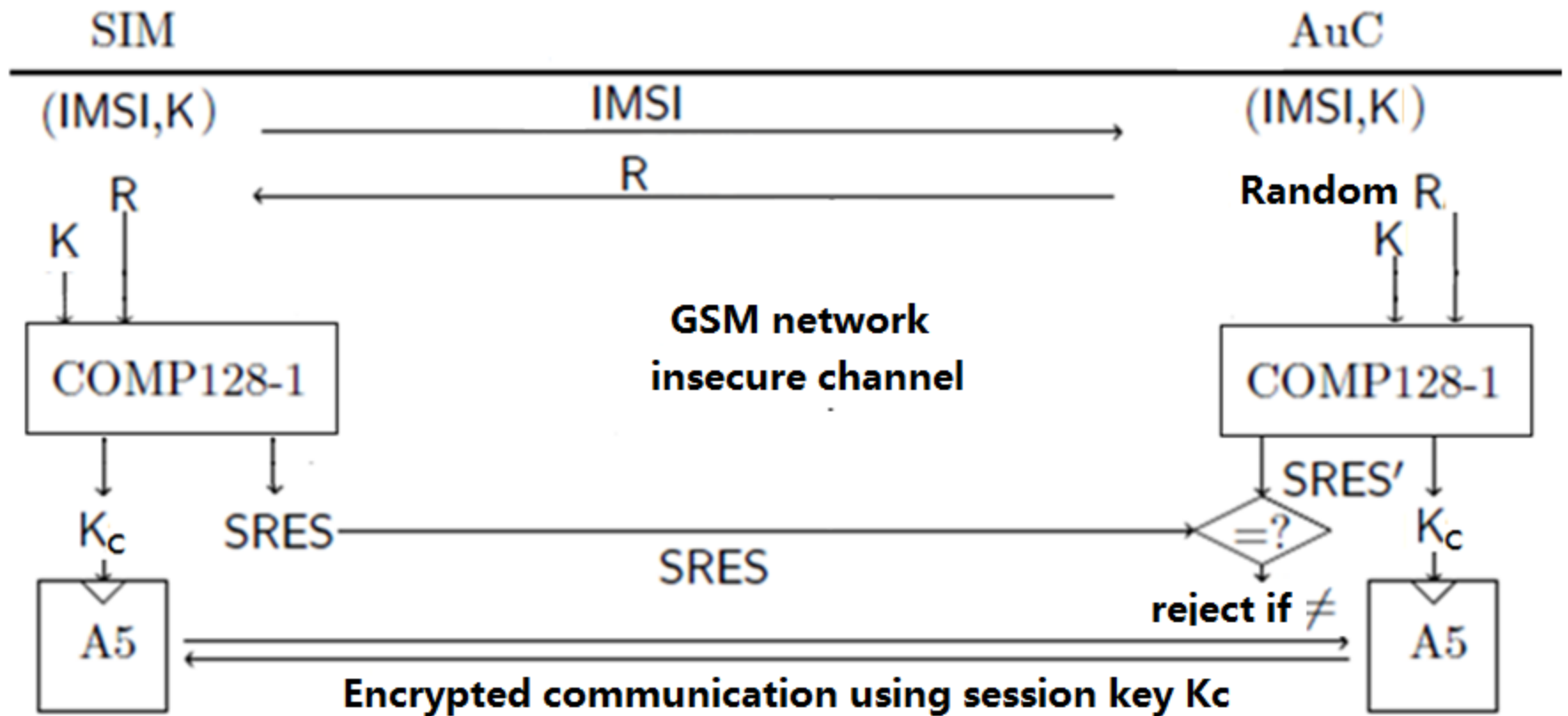


- The key of cloning a SIM card: recover the key K



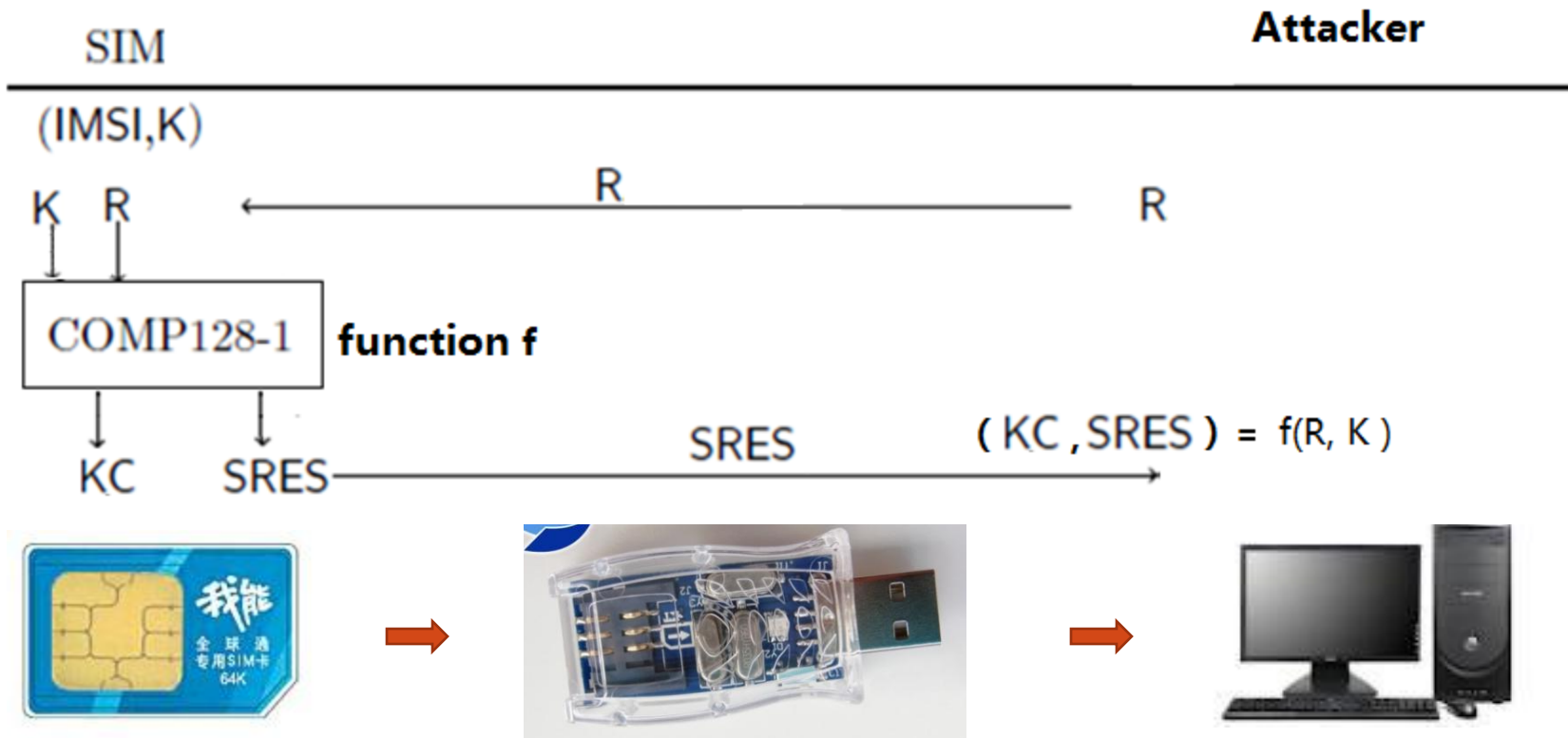
# Authentication between SIM card and base station (AuC)

GSM SIM uses the COMP128-1 algorithm for the authentication.



# Mathematical vs. physical attacks

- **Mathematical attack:** Attacker (impersonates the AuC), sends (**possibly malicious**) inputs  $R$  and observes output  $s$  accordingly, and try to recover  $K$ .



- **Side-channel attack:** In addition, attackers can capture some physical information such as power consumption.

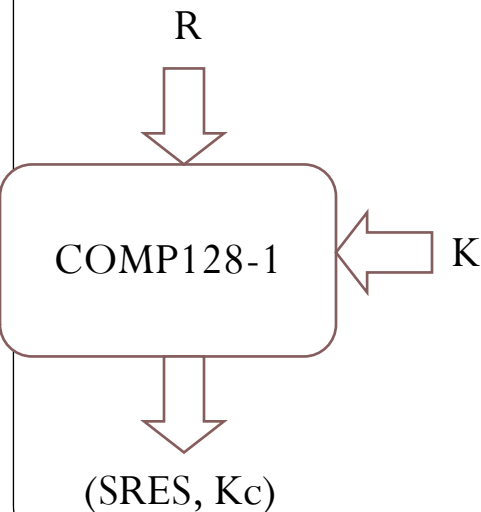
# History COMP128-1

- COMP128-1, as part of the GSM specification, drafted in 1987 and kept secret.
- In 1998, a research group at UC Berkeley (led by David Wagner) reversed engineered COMP128-1, and release it on the internet.
- COMP128-1 is a cryptographic hash function with a butterfly structure (FFT-HASH) .
- Targets of this work: a few SIMs cards from several (anonymized) manufacturers and operators.

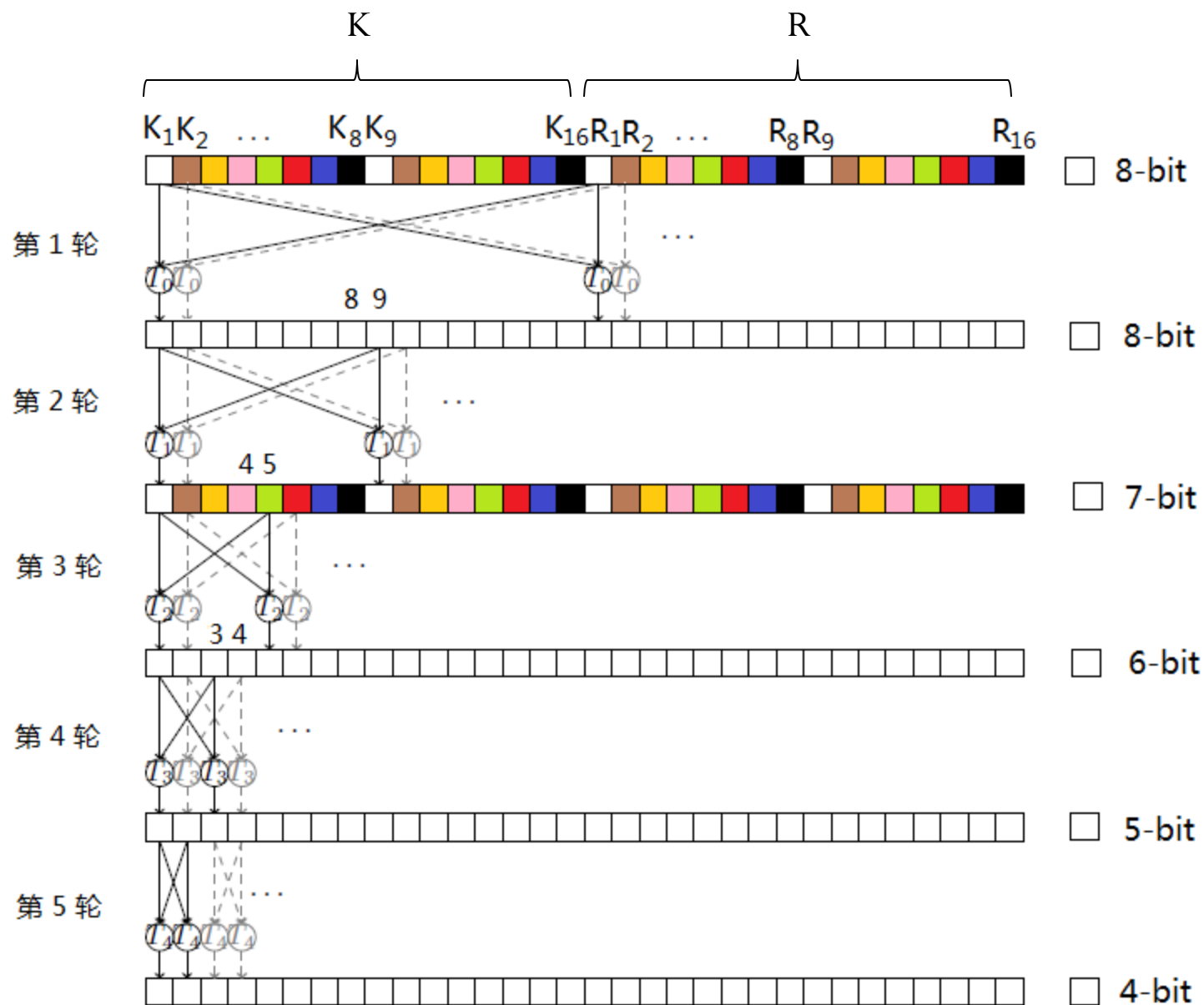
# Pseudo-code of COMP128-1

- COMP128-1 is cryptographic hash function.
- Input: 32-byte (i.e. 16-byte random R, 16-byte secret K)
- Output: 12-byte (i.e. 4-byte SRES 和 8-byte Kc).
- Pseudo-code:

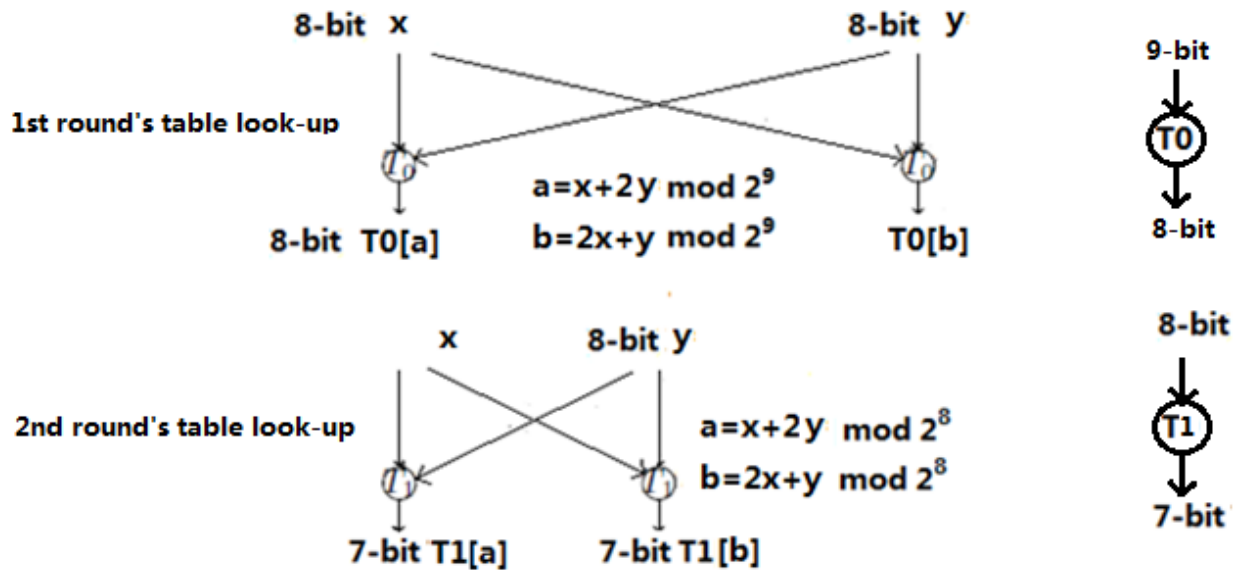
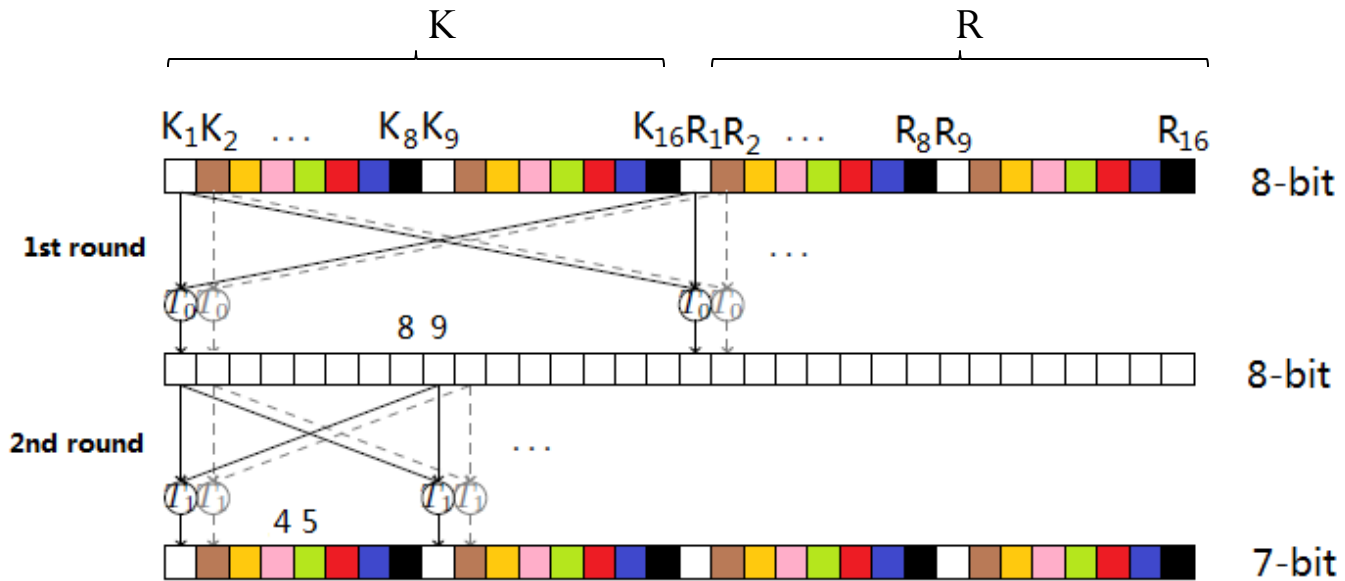
```
function COMP128-1(R, K)
begin
    for j=16 to 31 do           { * 调入随机数 R * }
        X[j] := R[j - 16];
    for i=0 to 7 do           { * 8次循环 * }
    begin
        for j=0 to 15 do      { * 调入密钥 K * }
            X[j] := K[j];
        call Compress;       { * 压缩函数 * }
        call FormBitsFromBytes; { * 格式转换 * }
        if i < 7 then        { * 置换 * }
            call Permute
        end;
    end;
end;
```



# Compression subroutine

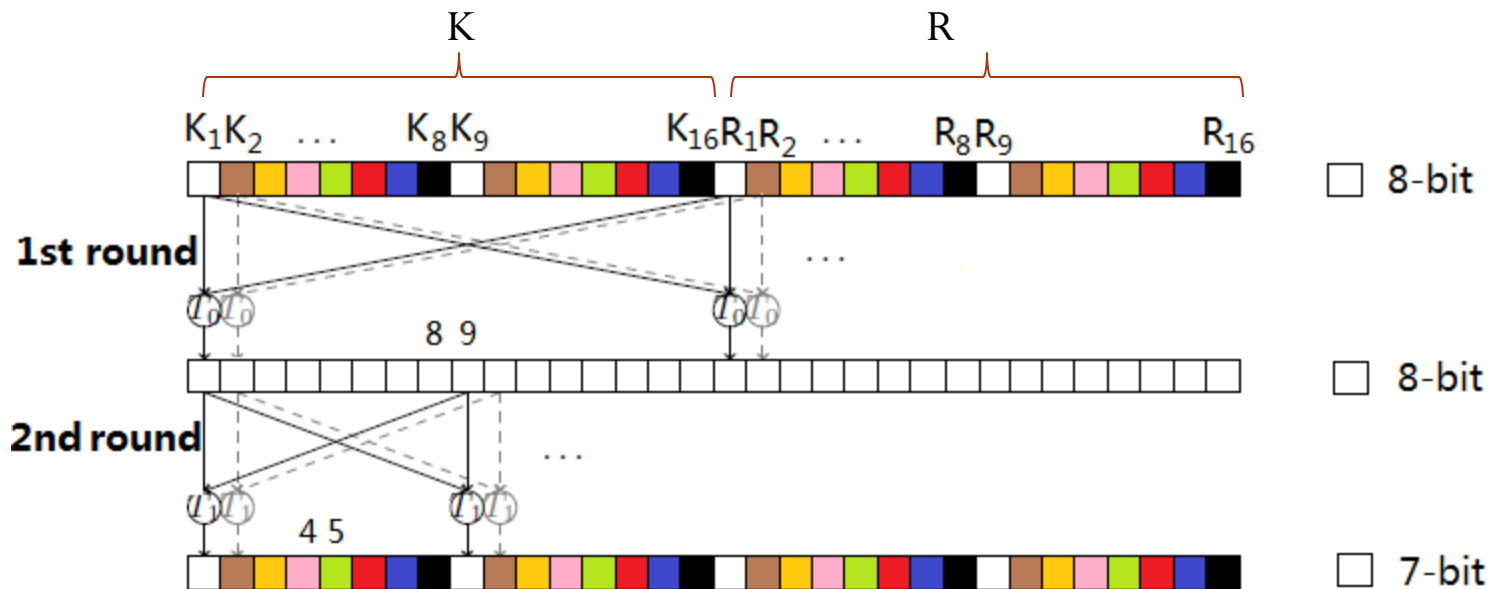


# flaw : insufficient diffusion



# Exploiting the Flaw: Collision attack

- Strategy: Divide and Conquer.
- Attack one color(1 key byte) at a time, fix the rest colors (s.t. collision on the output of 2<sup>nd</sup> round can propagate to the final output).
- Each color at 2<sup>nd</sup> round has 28 (4x7) bits, by birthday paradox, it takes  $2^{14}$  inputs to obtain 1 collision, so covering whole key needs  $2^{14} \times 8 = 131,000$  inputs.



# Collision attacks are implemented: SIM cloning kits available

- Low cost (~\$10).
- Cloning kit: SIM card reader, software (driver, cracking, SIM writing), blank SIM card
- Effective with COMP 128-1.



FREE  
shipping

## SIMMAX GSM 16-Number-in-1 SIM Card with USB Card Reader/Writer and Cloning Kit

Item condition: **New**

Time left: 1 day 9 hours (Apr 01, 2013 10:56:56 PDT)

Starting bid: **US \$9.99**

[ 0 bids ]

Enter US \$9.99 or more

Place bid

Price: **US \$14.99**

Buy It Now

Add to cart

Add to Watch list

Mouse over image to zoom

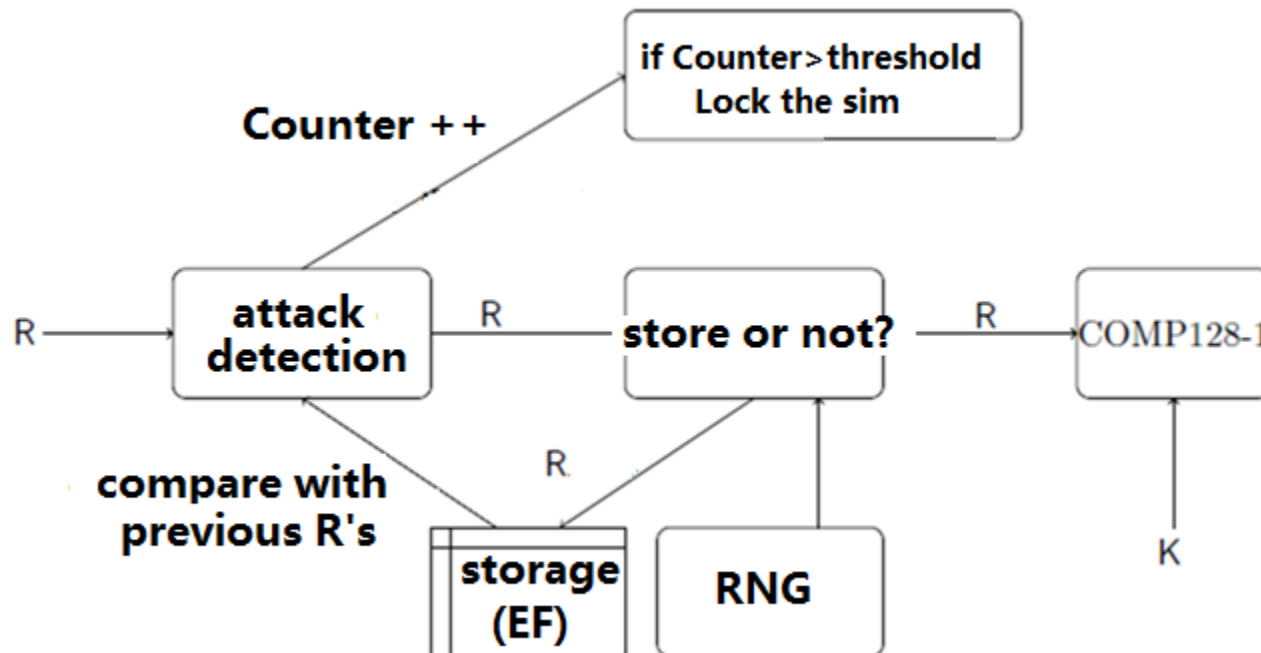


# Ad-hoc Countermeasures

- Move to newer versions COMP128-2, COMP128-3 (still kept secret!)
- Patch COMP128-1:

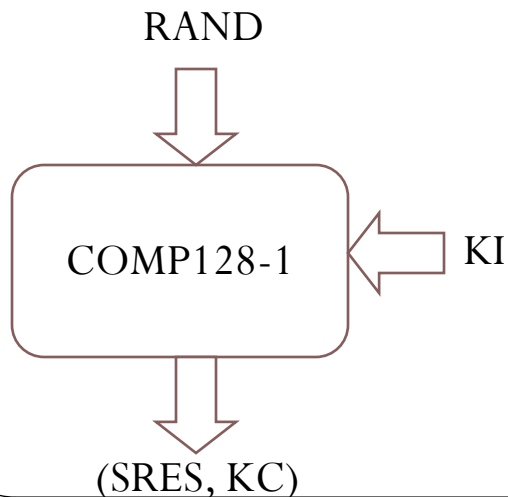
Known attacks easy to detect: attacker sends many correlated inputs.

Detecting heuristics (used by some operators): Store a few previous inputs, compare with the current one. Lock the card if too many attempts are detected.



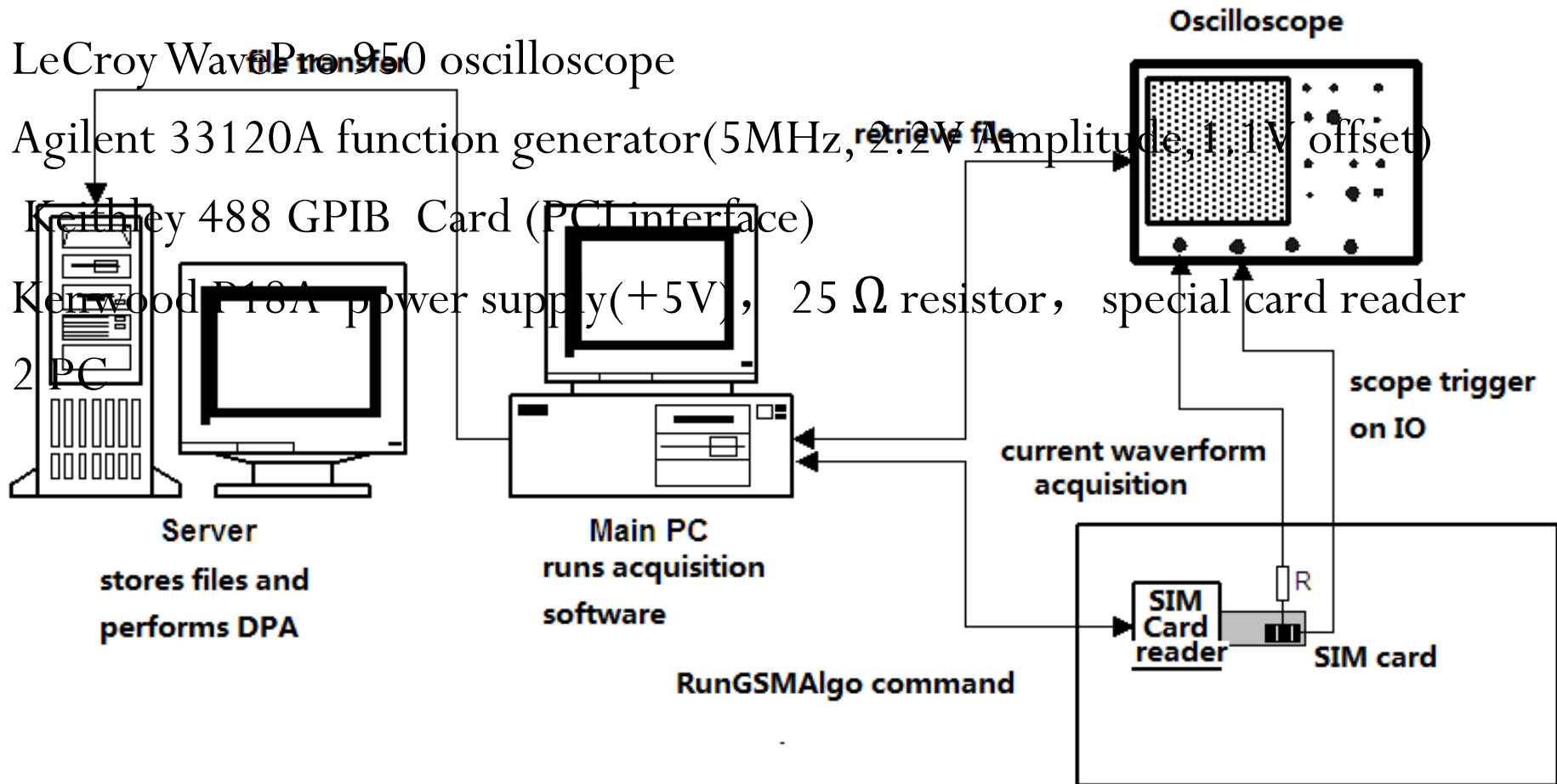
# Attack 2 (our results): Power Analysis Attacks

- Collision attacks fail because they are easy to detect.
- Power analysis: Send truly random R to SIM, not causing sim lock.
- How it works: SIM relies on external power and clocking signal.



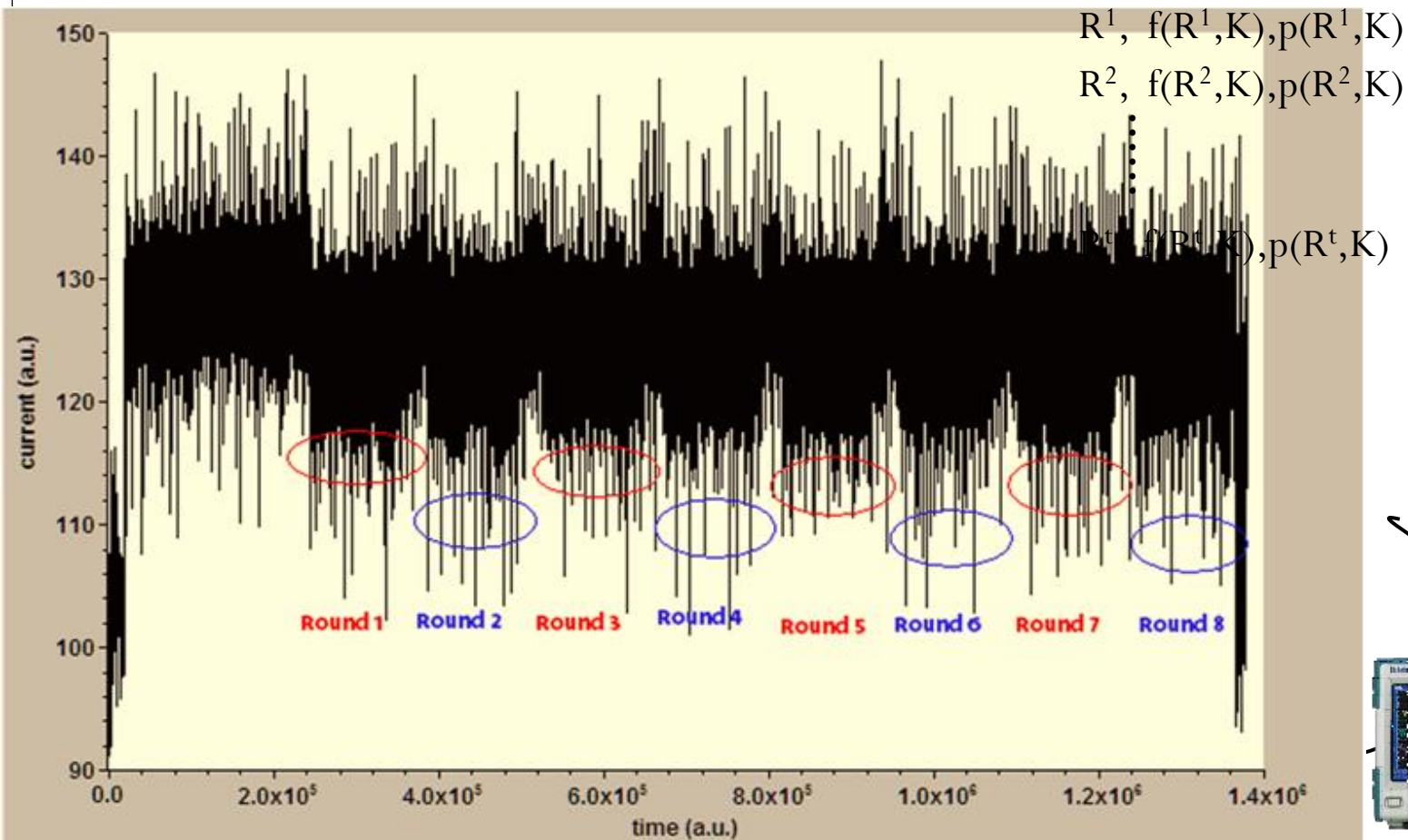
# Measurement Setup for Power Analysis

- LeCroy WavePro 950 oscilloscope
- Agilent 33120A function generator(5MHz, 2.2V Amplitude, 1.1V offset)
- Keithley 488 GPIB Card (PCI interface)
- Kenwood P18A power supply(+5V), 25  $\Omega$  resistor, special card reader
- 2 PC



# Power Trace Measurement

- Send random  $R$ , measure the corresponding output and power traces, and repeat.



Power trace  $p$

# How secrets are leaked from traces (leakage model)?

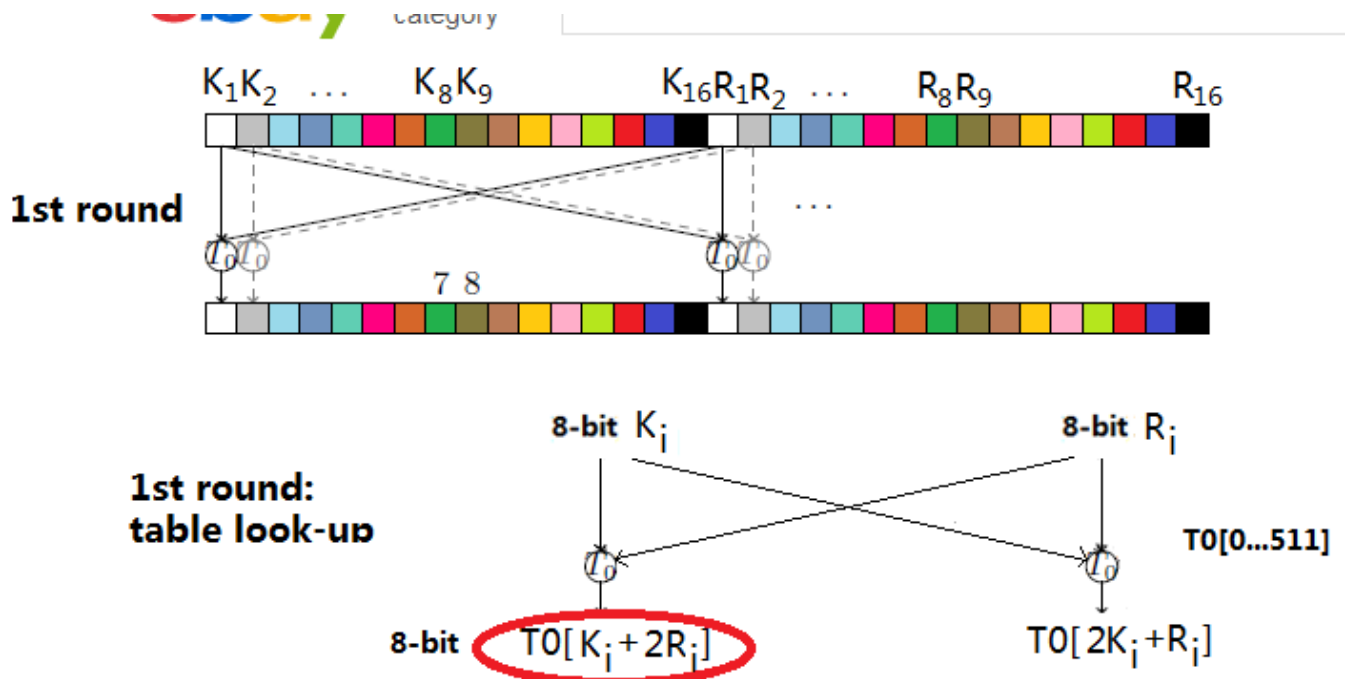
- Hamming weight model: The power consumption (for preserving value e.g.  $r=10100111$ ) is proportional (or conversely) to its Hamming weight.
- Applicable to CMOS circuits (with precharged data bus)

	time $t[i]$	time $t[i+1]$	Power ( $i \rightarrow i+1$ )
Byte[0]	0	1	$E_{0 \rightarrow 1}$
Byte[1]	0	0	$E_{0 \rightarrow 0}$
Byte[2]	0	1	$E_{0 \rightarrow 1}$
Byte[3]	0	0	$E_{0 \rightarrow 0}$
Byte[4]	0	0	$E_{0 \rightarrow 0}$
Byte[5]	0	1	$E_{0 \rightarrow 1}$
Byte[6]	0	1	$E_{0 \rightarrow 1}$
Byte[7]	0	1	$E_{0 \rightarrow 1}$

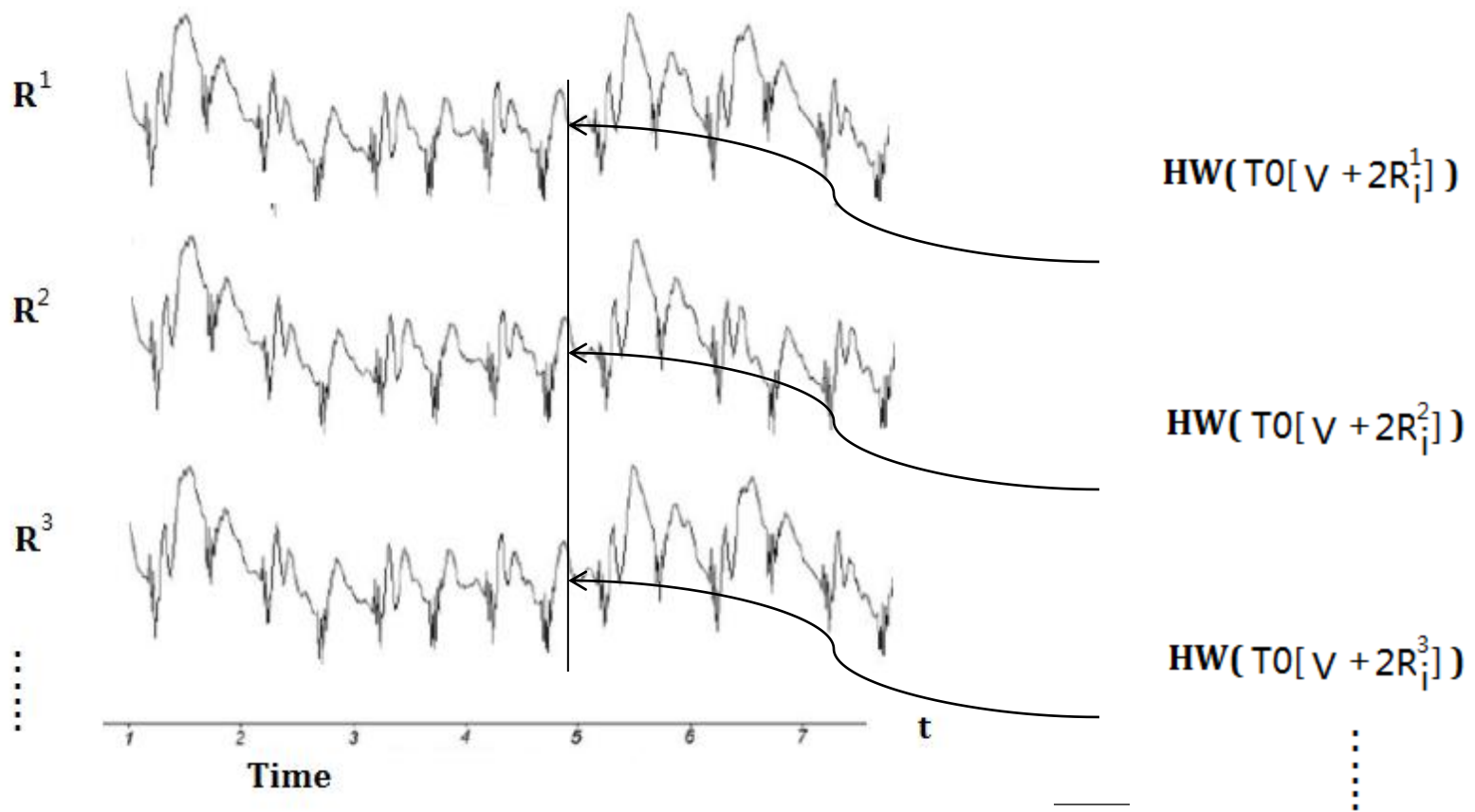
$$\text{Total: } 5E_{0 \rightarrow 1} + 3E_{0 \rightarrow 0} \approx 5E_{0 \rightarrow 1}$$

# Which intermediate result as the target?

- Strategy: Attack one color at a time ( $0 \leq i \leq 15$ ), but not fixing the rest colors (not causing SIM card lock).
- hypothesis testing: Target at  $T0[K_i + 2R_i]$ , assume  $K_i = v$  (256 possibilities), compute the correlation coefficient between  $T0[v + 2R_i]$ 's Hamming weight and power traces.
- For correct guess  $K_i = v$ , the correlation should be maximal.

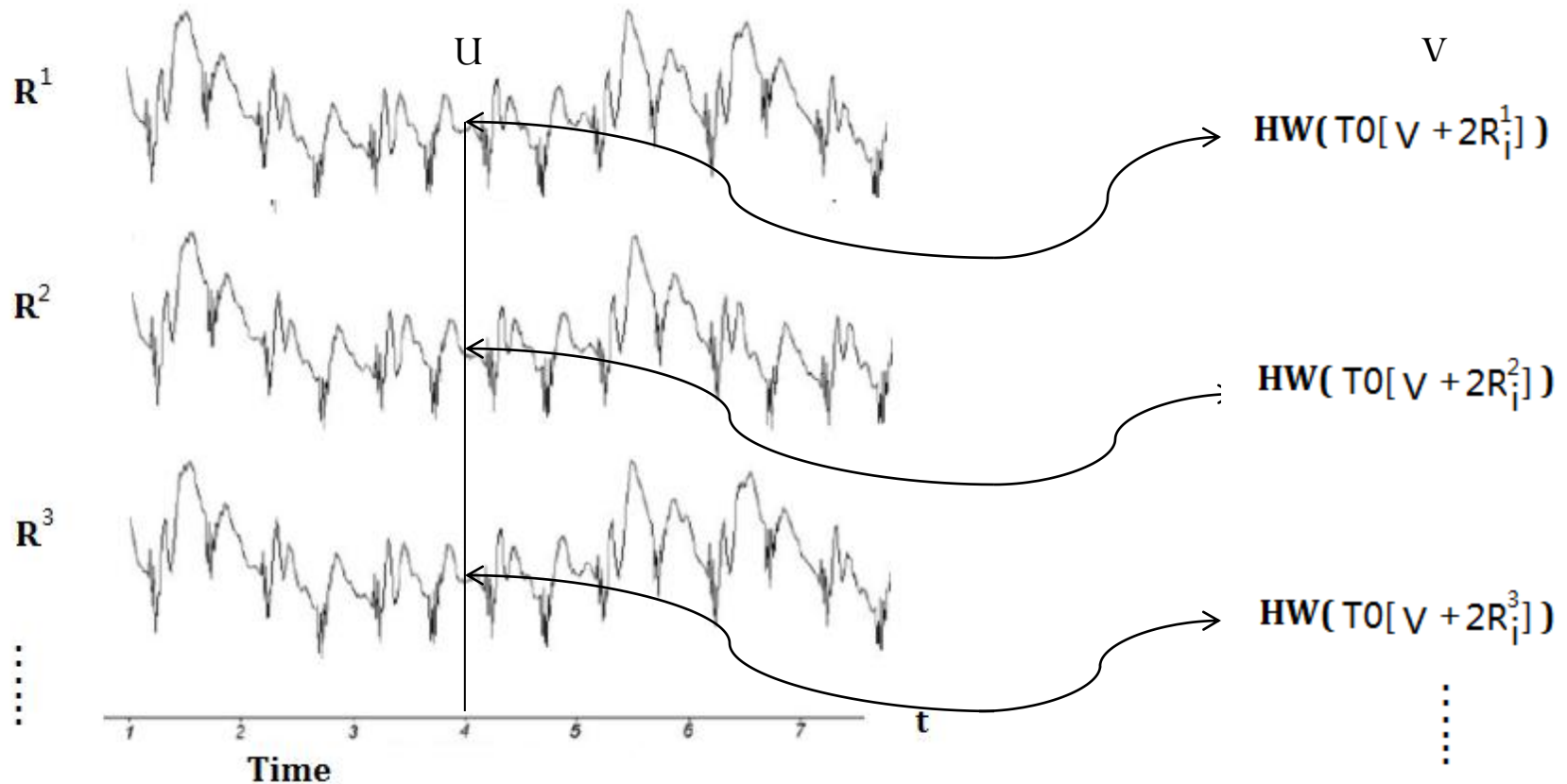


# Traces might be misaligned



Assume  $K_i = v$ , Compute correlation coefficient ( between power traces and  $HW(T_0[v + 2R_i])$  )

- hypothesis testing: compute the coefficient corresponding to  $v=0,1,\dots,255$  one by one, the maximum should be with the correct hypothesis.





# Pearson correlation coefficient

Correlation coefficient between  $U$  and  $V$ , denoted by  $\rho_{U,V}$ , is:

$$\rho_{U,V} \stackrel{\text{def}}{=} \frac{\mathbf{E}[(X - \mu_U)(Y - \mu_V)]}{\sigma_U \sigma_V}$$

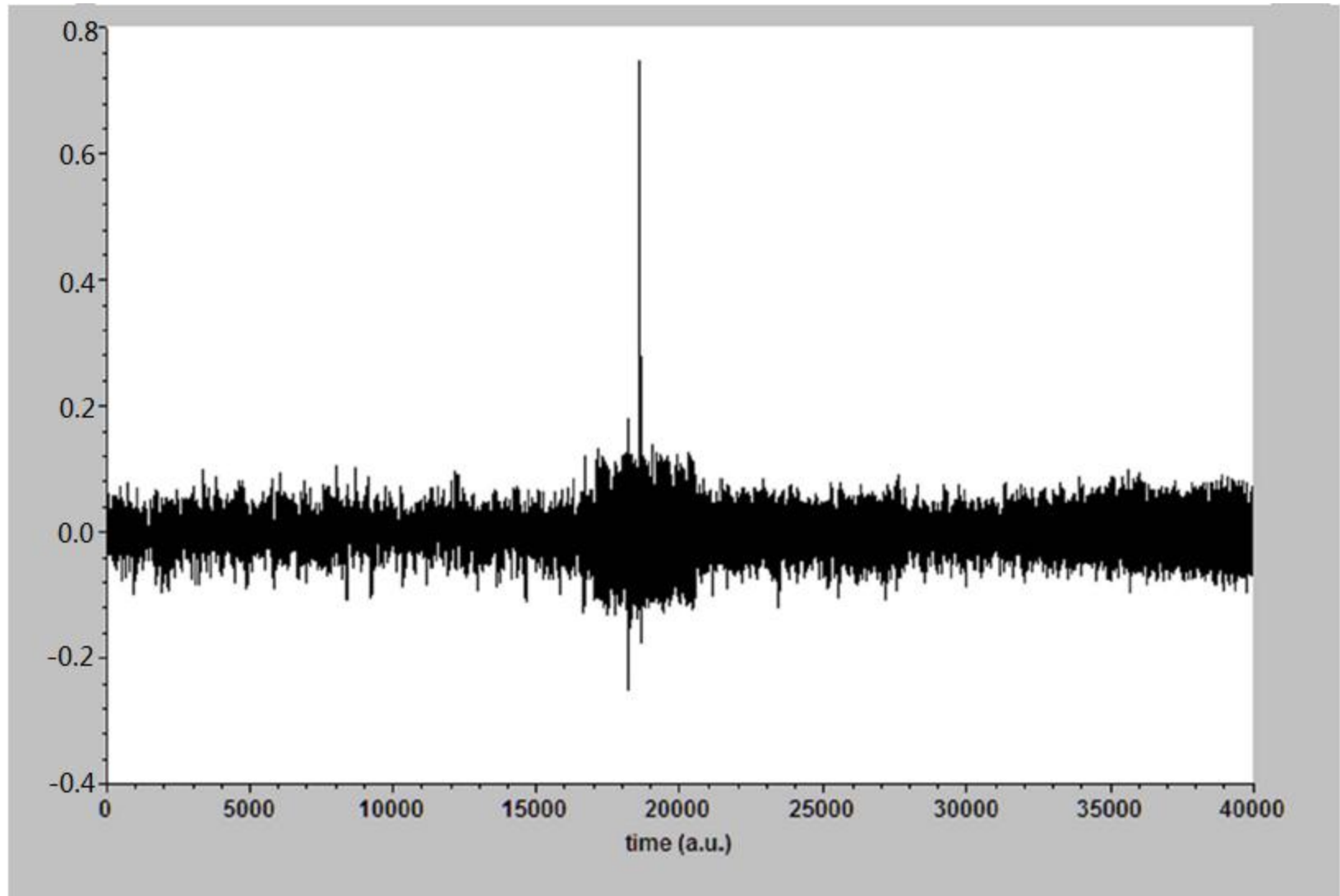
where  $\mathbf{E}$  is expectation,  $\mu_U \stackrel{\text{def}}{=} \mathbf{E}[U]$ , and standard deviation  $\sigma_U \stackrel{\text{def}}{=} \sqrt{\mathbf{E}[(U - \mu_U)^2]}$ .

By sampling from  $(U, V)$  to  $(u_1, v_1), (u_2, v_2), \dots, (u_n, v_n)$ , the estimator of  $\rho_{X,Y}$ , denoted by  $r_{x,y}$ , is given by:

$$r_{x,y} = \frac{\sum_{i=1}^n (u_i - \bar{u})(v_i - \bar{v})}{\sqrt{\sum_{i=1}^n (u_i - \bar{u})^2} \sqrt{\sum_{i=1}^n (v_i - \bar{v})^2}},$$

where  $\bar{u} = \frac{u_1 + u_2 + \dots + u_n}{n}$  and  $\bar{v} = \frac{v_1 + v_2 + \dots + v_n}{n}$  denotes mean value.

coefficient for a correct hypothesis ( $K_i=v$ )



# Power analysis vs. collision attacks

- Targets: 4 SIM cards from two mobile operators and 4 different manufacturers
- Efforts in terms of: the number of inputs (traces) needed.

	manufacturer	operator	patch (countermeasure)	DPA	collision attacks
SIM#1	I	A	Not available	400	20,000
SIM#2	II	B	I-C	200	$\geq 20,000$
SIM#3	III	B	I-C + C-F	4000	fail (card locked)
SIM#4	IV	B	I-C + C-F	10000	fail (card locked)

- Collision attacks: cheap set-up, only applicable to unpatched targets
- Power analysis: powerful, provided with special measurement setup

# Lessons Learned

- Awareness of physical security for small embedded devices.
- The contrast:
  - Low cost devices  $\approx$  limited budget for CC/EMVCo security testing.
  - Low-cost  $\times$  huge volume = big impact / loss
- Some SIM cards are used for more sensitive applications such as mobile payments.
- Practical security requires BOTH:
  - A mathematically secure (and publicly referred) algorithm.
  - Sufficient countermeasures in place against physical attacks.

Thanks!

